

Crypto Engineering (GBX9SY03)

TP — Generic second preimage attacks on long messages for narrow-pipe Merkle-Damgård hash functions

2019-01-11

Grading

This TP is graded as part of the *contrôle continu*. You must send a written report (in a portable format) detailing your answers to the questions, and the corresponding source code, *including all tests*, **with compilation and execution instructions** by the end of the second session, (2019-01-16T23:59+0200) to:

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Working in teams of two is allowed and encouraged (but not mandatory), in which case only one report needs to be sent, with the name of both students clearly mentioned.

Introduction

The goal of this TP is to implement a generic second preimage attack for long messages for Merkle-Damgård hash functions, in the specific case where the compression function used within the hash function follows a Davies-Meyer construction. This is a particular case of the attacks (re-)designed by Kelsey and Schneier [KS05].

You will have to implement and run the full attack on a toy hash function with 48-bit hashes based on the SPECK48/96 block cipher [BSS⁺13]. While this hash function will obviously be vulnerable to brute-force attacks (and maybe even more), it is simple to implement and it reasonably behaves like a random function, which is all we need here.

Part one: preparatory work

Download the specifications of SPECK48/96 at <http://eprint.iacr.org/2013/404> and the tarball at https://www-ljk.imag.fr/membres/Pierre.Karpman/cry_eng2018_tp_second_preim.tar.bz2. This tarball contains the file `second_preim_48.c` that already partially implements some functions, which you will be required to complete.

Question 1

Finish to implement the function `speck48_96` of the encryption with SPECK48/96, and verify the correctness of your implementation with the test values provided in [BSS⁺13, App. C]. Write for that purpose a function `int test_sp48(void)`. (Be careful about the order in which the 24-bit words of the key and message are provided to your function.)

[https:](https://www-ljk.imag.fr/membres/Pierre.Karpman/cry_eng2018_tp_second_preim.pdf)

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Question 2

Implement the function `speck48_96_inv` of the decryption by SPECK48/96. Write a test function `int test_sp48_inv(void)` that verifies that `speck48_96_inv` and `speck48_96` are inverses of each other.

Question 3

Implement the function `cs48_dm` that transforms SPECK48/96 into a compression function using a Davies-Meyer construction with an XOR feedforward, i.e. defining $\mathcal{F}(h, m)$ as $\mathcal{E}(m, h) \oplus h$. Write a test function `int test_cs48_dm(void)` to check that the result of calling the compression function on the all-zero input is `0x7FDD5A6EB248ULL`.

Question 4

Implement the function `get_cs48_dm_fp` that returns the unique fixed-point `fp` for the function `cs48_dm` with a fixed message `m`; that is, compute the unique 48-bit string `fp` such that `cs48_dm(m, fp) == fp`. Write a test function `int test_cs48_dm_fp(void)`.

Part two: the attack

The main idea of the second-preimage attack that we consider here is to search for collisions with one of the intermediate chaining values of a long message.

Given an ℓ -block message $M = m_1 || \dots || m_\ell$, the computation of $\mathcal{H}(m)$ with a narrow-pipe Merkle-Damgård hash function \mathcal{H} with compression function \mathcal{F} requires the computation of ℓ chaining values $h_i = \mathcal{F}(h_{i-1}, m_i)$. Assuming first a simplistic function \mathcal{H} that does not use any sort of padding, one can see that an attacker who has found M' s.t. $\mathcal{H}(M') = h_i$ for some $0 \leq i \leq \ell$ may form the message $M'' = M' || m_{i+1} || \dots || m_\ell$ leading to h_ℓ , and thus found a second preimage M'' for M . The expected number of tries for finding a collision between an n -bit chaining value with ℓ targets being $2^{n-\log(\ell)}$, this attack is better than a generic search, and has minimum complexity $2^{n/2}$.

In reality, \mathcal{H} will use some form of padding that includes the length of the message being hashed. If M and its potential second preimage M'' are not of the same length, their padding (added after their last message block and before the computation of the hash) will be distinct and thus lead to different hashes with overwhelming probability, thus invalidating the attack.

One way to go around this problem is to first compute an *expandable message* E that will allow to “lengthen” M'' up to the boundary of the block i while preserving the collision with h_i . One may then build $M''' = E || M''$ which is of the same length as M and thus has the same padding, resulting in a true second-preimage.

You are required to implement this attack for the hash function `hs48`, already implemented in the provided file. This function takes messages made of an integral number of 96-bit *blocks*, each represented as a 4-double-word array `uint32_t b[4]` where only the three lowest bytes of each double-word is set.

Question 1

Implement the function `find_exp_mess`. This function must return two one-block messages `m1` and `m2` such that there exists a value `h` equal to both `cs48_dm(m1, IV)` and

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`get_cs48_dm_fp(m2)`. In other words, for any message `m` made of one copy of `m1` followed by `n >= 1` copies of `m2`, one has `hs48(m,n+1,0,0) == h`. Write a test function `int test_em(void)` to validate your function.

Give a detailed description of your implementation of `find_exp_mess` (for instance describing your choices for the data structures) and evaluate its performance (both theoretically and practically).

How to proceed. To answer this question, you are advised to perform a meet-in-the-middle search: compute N possible chaining values for N random first-block messages `m1`, and then compute fixed-points for random messages until one of them collides with a known chaining value. You may use the pseudo-random number generator provided in the header file `xoshiro256starstar.h`.*

With a reasonably good implementation, finding an expandable message should take significantly less than one minute on average.

Question 2

Implement the function `attack`. This function must return a second preimage for the message `mess` of 2^{18} blocks generated by the following code:

```
for (int i = 0; i < (1 << 20); i+=4)
{
    mess[i + 0] = i;
    mess[i + 1] = 0;
    mess[i + 2] = 0;
    mess[i + 3] = 0;
}
```

whose hash is equal to `0xFF3FD9D23B89`.

Give a detailed description of your implementation and of its performance (both theoretical and practical). Be careful to document enough the output of this function so that it provides all the information needed to characterise the second preimage (you may for instance freely use the “verbose” option of `hs48`).

How to proceed. You should first compute an expandable message with associated fixed-point `fp`, and then search for a collision block `cm` s.t. `cs48_dm(cm,fp)` is equal to one of the chaining values of `mess`. Once such a block is found, you need to form the second preimage message `mess2` by expanding the expandable message to an appropriate number of blocks and suffixing the colliding block and the remaining blocks identical to the ones of `mess`. Finally compute the hash of `mess2` to validate your attack.

With a reasonably good implementation, finding a full attack should take significantly less than ten minutes on average.

References

- [BSS⁺13] Ray Beaulieu, Douglas Shors, Jason Smith, Stefan Treatman-Clark, Bryan Weeks, and Louis Wingers, *The SIMON and SPECK Families of Lightweight Block Ciphers*, IACR Cryptology ePrint Archive **2013** (2013), 404.

*See also <http://xoshiro.di.unimi.it/>.

[https:](https://www-ljk.imag.fr/membres/Pierre.Karpman/cry_eng2018_tp_second_preim.pdf)

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- [KS05] John Kelsey and Bruce Schneier, *Second Preimages on n -Bit Hash Functions for Much Less than 2^n Work*, Advances in Cryptology — EUROCRYPT 2005 (Ronald Cramer, ed.), Lecture Notes in Computer Science, vol. 3494, Springer, 2005, pp. 474–490.