



## Symmetric Determinantal Representations of Polynomials in Characteristic 2

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# Introduction

$$xy + yz + xz$$



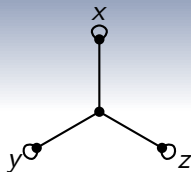
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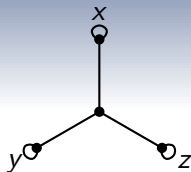
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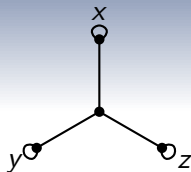
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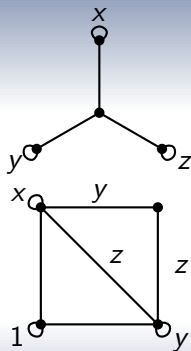




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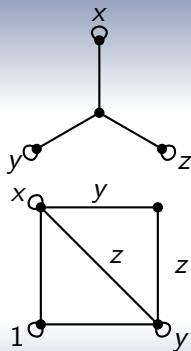




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What about  $xy + z$ ?





# Formalization

## Symmetric Determinantal Representation

SDR of  $p \in \mathbb{F}[x_1, \dots, x_m]$ :

- **Symmetric** matrix  $M$ ;
- Entries: elements of  $\mathbb{F} \cup \{x_1, \dots, x_m\}$ ;
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A polynomial is said **representable** if it has a SDR.



# Problems

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Are **all** polynomials representable in characteristic 2?



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- 2 Polytime algorithm for **deciding representability**.
- 3 Polytime algorithm for **finding SDRs**.



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# Motivation

- Algebraic Complexity



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  - Universality of the determinant  
[Valiant'79, Toda'92, Malod & Portier'06]





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  - SDRs in characteristic  $\neq 2$  always exist  
[G., Kaltofen, Koiran, Portier'11]
- Convex Optimization



# Determinant and cycle covers

## Determinant

$\mathfrak{S}_n =$  Permutation group of  $\{1, \dots, n\}$

$$\det A = \sum_{\sigma \in \mathfrak{S}_n} (-1)^{\text{sgn}(\sigma)} \prod_{i=1}^n A_{i, \sigma(i)}$$



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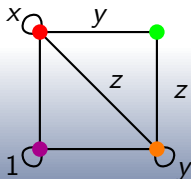
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●				
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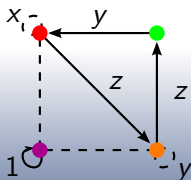
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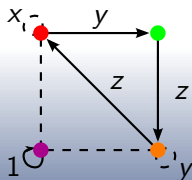
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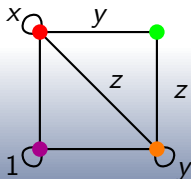
# Determinant and partial matchings

## Determinant in characteristic 2 of symmetric matrices

$\mathfrak{I}_n =$  Involutions of  $\{1, \dots, n\}$

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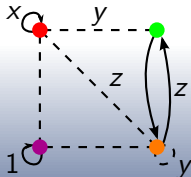
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# Representable polynomials



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### *Lemma*

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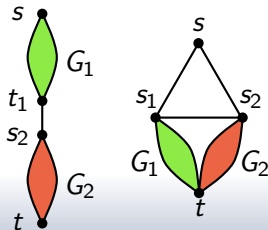
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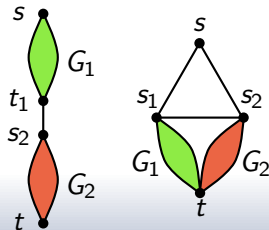
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## Lemma

For all  $P$ ,  $P^2$  is representable.

- $\det(G \setminus \{s, t\}) = 1$
- $\det(G \setminus \{s\}) = \det(G \setminus \{t\}) = 0$





## A class of representable polynomials

### *Theorem*

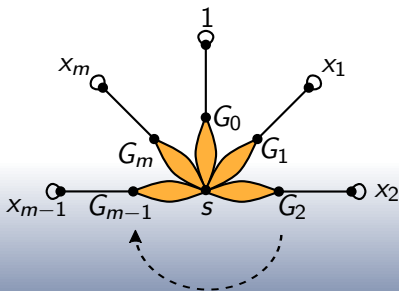
$L(x_1, \dots, x_m) = P_0^2 + x_1 P_1^2 + \dots + x_m P_m^2$  is representable.



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## Obstructions to representability

### *Theorem*

*If  $P$  is representable, then*

$$P \equiv L_1 \times \cdots \times L_k \pmod{\langle x_1^2 + 1, \dots, x_m^2 + 1 \rangle}$$

*where the  $L_i$ 's are linear.*

*(linear = degree-1)*



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*where the  $L_i$ 's are linear, and the  $\ell_i$ 's are squares.*

$P$  is said **factorizable** modulo  $\langle x_1^2 + \ell_1, \dots, x_m^2 + \ell_m \rangle$ .



## Proof idea

- Modulo  $\langle x_1^2 + \ell_1, \dots, x_m^2 + \ell_m \rangle$  : no variable **outside the diagonal**

$$xz + y^2 = \det \begin{pmatrix} x & y \\ y & z \end{pmatrix}$$



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$$xz + y^2 = \det \begin{pmatrix} x & y \\ y & z \end{pmatrix} \equiv \det \begin{pmatrix} x & 1 \\ 1 & z \end{pmatrix} \pmod{\langle x^2 + 1, y^2 + 1, z^2 + 1 \rangle}$$



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$$xz + y^2 = \det \begin{pmatrix} x & y \\ y & z \end{pmatrix} \equiv \det \begin{pmatrix} x & 1 \\ 1 & z \end{pmatrix} \equiv \det \begin{pmatrix} x & 1+x \\ 1+x & x+z \end{pmatrix} \\ \text{mod } \langle x^2 + 1, y^2 + 1, z^2 + 1 \rangle$$



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# Multilinear polynomials

## *Theorem*

Let  $P$  be a **multilinear** polynomial. The three following propositions are equivalent:

- (i)  $P$  is representable.
- (ii)  $\exists \ell, P$  is factorizable modulo  $\langle x_1^2 + \ell_1, \dots, x_m^2 + \ell_m \rangle$ .
- (iii)  $\forall \ell, P$  is factorizable modulo  $\langle x_1^2 + \ell_1, \dots, x_m^2 + \ell_m \rangle$ .



## Proof idea

$$\textcircled{1} \quad xy + yz + xz \equiv xyz(x + y + z) \pmod{\langle x^2 + 1, y^2 + 1, z^2 + 1 \rangle}$$



## Proof idea

①  $xy + yz + xz \equiv xyz(x + y + z) \pmod{\langle x^2 + 1, y^2 + 1, z^2 + 1 \rangle}$

②  $xyz(x + y + z) = \det \begin{bmatrix} x & 0 & 0 & & & \\ 0 & y & 0 & & & \\ 0 & 0 & z & & & \\ & & & x & 1 & 0 & 0 & 0 \\ & & & 1 & 0 & 1 & 0 & 0 \\ & & & 0 & 1 & y & 1 & 0 \\ & & & 0 & 0 & 1 & 0 & 1 \\ & & & 0 & 0 & 0 & 1 & z \end{bmatrix}$



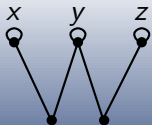


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②  $xyz(x + y + z) \equiv \det$

$$\begin{bmatrix} x & 0 & 0 & & & \\ 0 & y & 0 & & & \\ 0 & 0 & z & & & \\ & & & x & 1 & 0 & 0 & 0 \\ & & & 1 & 0 & 1 & 0 & 0 \\ & & & 0 & 1 & y & 1 & 0 \\ & & & 0 & 0 & 1 & 0 & 1 \\ & & & 0 & 0 & 0 & 1 & z \end{bmatrix}$$







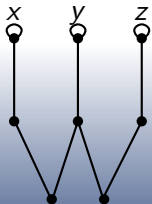


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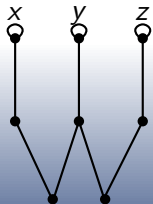


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$$\begin{bmatrix} x & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & y & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & z & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} = xy + yz + xz$$





## Non-multilinear polynomials

$$xyz^2 + y^3z + x^3y^2z \longrightarrow \xi_z^2 \cdot xy + \xi_y^2 \cdot yz + \xi_x^2 \xi_y^2 \cdot xz$$



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### *Theorem*

- *Factorizable in  $\mathbb{F}[\xi_x, \xi_y, \xi_z][x, y, z] \implies$  representable*
- ***Not** factorizable in  $\mathbb{F}(\xi_x, \xi_y, \xi_z)[x, y, z] \implies$  not representable*



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## Conclusion

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- Characterization for multilinear polynomials
- Sufficient and necessary conditions for other polynomials



## Conclusion

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- Characterization for multilinear polynomials
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Wait! Is  $xy + z$  representable?



## Finding a factor

$$(x + y + z + 1) \times (x + y + z + 1) \times \cdots \times (x + y + z + 1) \\ \stackrel{?}{\equiv} xy + z \pmod{\langle x^2, y^2, z^2 \rangle}$$



## Finding a factor

$$(x + y + z + 1) \times (x + y + z + \mathbf{1}) \times \cdots \times (x + y + z + \mathbf{1})$$
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## Finding a factor

$$\begin{aligned} & ( \quad z \quad ) \times (x + y + z + \mathbf{1}) \times \cdots \times (x + y + z + \mathbf{1}) \\ & \stackrel{?}{\equiv} xy + z \pmod{\langle x^2, y^2, z^2 \rangle} \end{aligned}$$



## Finding a factor

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- $\text{Lin}(xy + yz + y + z + 1) = y + z + 1$
- $\mathcal{I}_0 = \langle x^2, y^2, z^2 \rangle$



## Finding a factor

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### *Lemma*

Let  $P$  s.t.  $P(0) = 0$ ,  $\text{Lin}(P) \neq 0$ . Then

$$P \equiv L_1 \cdots L_k \pmod{\mathcal{I}_0} \implies L_i = \text{Lin}(P) \text{ for some } i.$$



# Divisibility

## *Lemma*

- $P$ : *multilinear polynomial*,
- $L$ : *monic linear polynomial s.t.  $L(0) = 0$ .*

*Then*

$$\exists Q, P \equiv L \times Q \pmod{\mathcal{I}_0} \implies P \equiv L \times \frac{\partial P}{\partial x} \pmod{\mathcal{I}_0}.$$



## Preparation

### *Lemma*

*Let  $P$  be a multilinear polynomial. Then there exists  $Q$  and a multilinear  $P^*$  such that*

- $P^* \equiv P \times Q \pmod{\langle x^2 + 1, y^2 + 1, z^2 + 1 \rangle}$
- $\text{Lin}(P^*) \neq 0$
- $P^*(0) = 0$



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*Let  $P$  be a multilinear polynomial. Then there exists  $Q$  and a multilinear  $P^*$  such that*

- $P^* \equiv P \times Q \pmod{\langle x^2 + 1, y^2 + 1, z^2 + 1 \rangle}$
- $\text{Lin}(P^*) \neq 0$
- $P^*(0) = 0$

$\rightsquigarrow$   $P^*$  is representable if and only if  $P$  is representable.



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**Example.** Suppose  $P = xy + yz + x + 1$ . Then  $P \times y \equiv x + z + xy + y \pmod{\langle x^2 + 1, y^2 + 1, z^2 + 1 \rangle}$ .



## An algorithm

ISREPRESENTABLE( $P$ ):

$P$  multilinear

① If  $P$  is linear Then RETURN TRUE





# An algorithm

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- 2 Compute  $P^*$ , multilinear, and  $Q$  such that
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$\rightsquigarrow \text{MERGE}_\ell(M_1, M_2)$ : SDR of  $Q$  (linear time)





## New algorithm

SDR( $P$ ):

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  - Then
    - $M_0 \leftarrow \text{SDR}(P_0)$
    - $M_1 \leftarrow \text{MAT}(\text{Lin}(P^*))$
    - $M_2 \leftarrow \text{MAT}(Q)$
    - RETURN MERGE<sub>1</sub>( $M_2, \text{MERGE}_0(M_1, M_0)$ )
  - Else RETURN FALSE



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# Summary

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*Thank you!*