Symmetric Determinantal Representations of Polynomials

Bruno Grenet[*†]

Joint work with Erich L. Kaltofen[‡], Pascal Koiran[*†] and Natacha Portier[*†]

[*]MC2 – LIP, ÉNS Lyon
[†]Theory Group – DCS, U. of Toronto
[‡]Dept. of Mathematics – North Carolina State U.

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The problem

$$(x + 3y)z = \det \begin{pmatrix} 0 & x & 3 & 0 & 0 \\ 0 & 1 & 0 & 0 & z \\ 0 & 0 & 1 & y & 0 \\ 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{pmatrix}$$

- Formal polynomial

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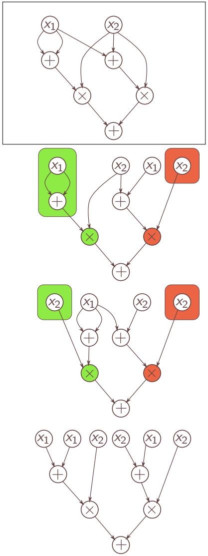
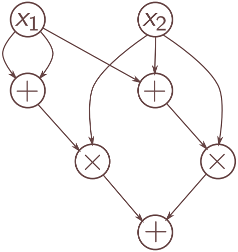
- Formal polynomial
- Smallest possible dimension of the matrix

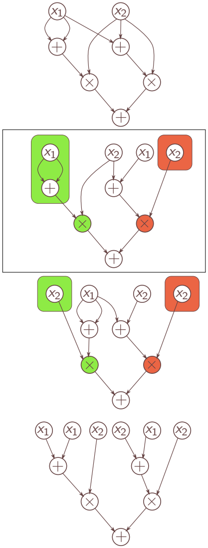
Representations of polynomials

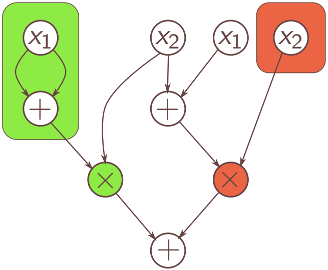




Circuit

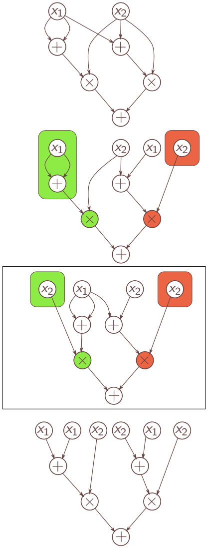
Representations of polynomials

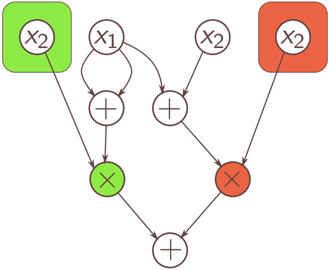




Weakly-skew circuit

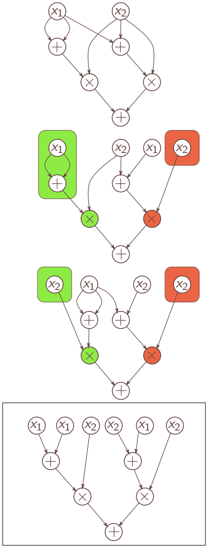
Representations of polynomials

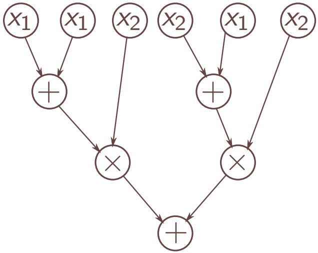




Skew circuit

Representations of polynomials





Formula

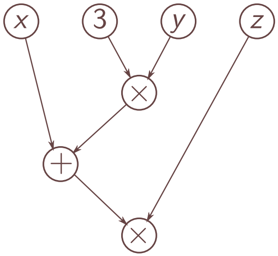
Motivation

L. G. Valiant, Completeness classes in algebra, STOC 79
$\rightsquigarrow$ Universality of the determinant

"We conclude that for the problem of finding a subexponential formula for a polynomial when one exists, linear algebra is essentially the only technique in the sense that it is always applicable".

An example

$(x + 3y)z$



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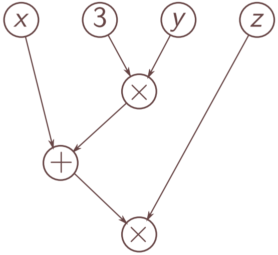

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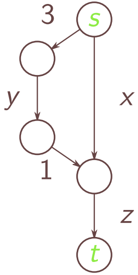
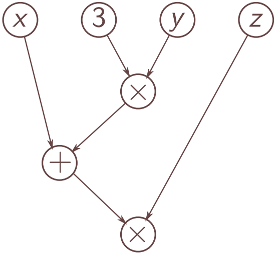
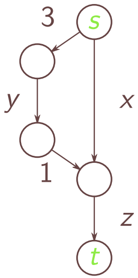




Arithmetic Branching Program

An example



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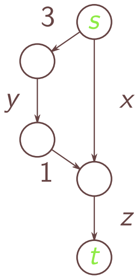

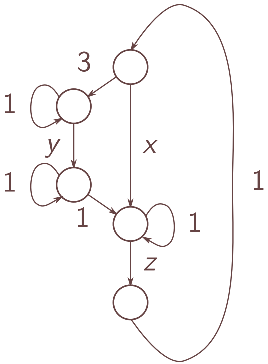

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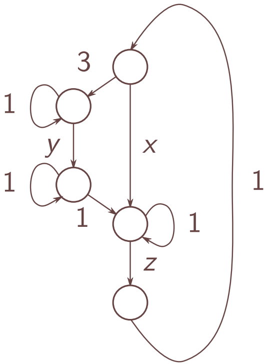

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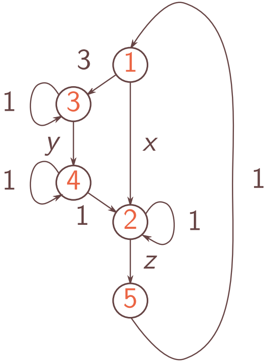

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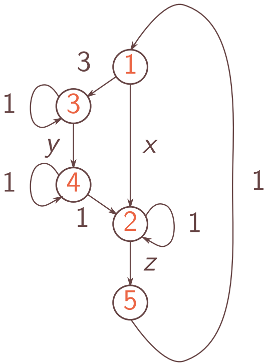

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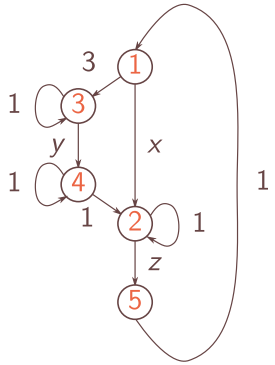

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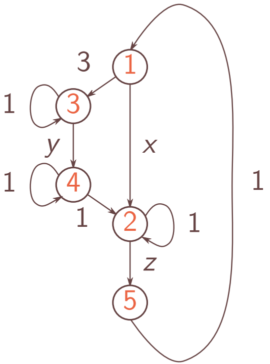

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- Up to signs, $\det A$ = sum of the weights of cycle covers in G

Outline

1. From polynomials to determinants
2. From polynomials to determinants of symmetric matrices
3. Characteristic 2

Upper bounds

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- $e+1$ if there is at least one addition in the formula: H. Liu and K.W. Regan, in *Improved construction for universality of determinant and permanent* (Inf. Process. Lett., 2006)

Valiant's construction (1/3)

- Input: a formula representing a polynomial $\varphi \in \mathbb{K}[X_1, \ldots, X_n]$ of size e

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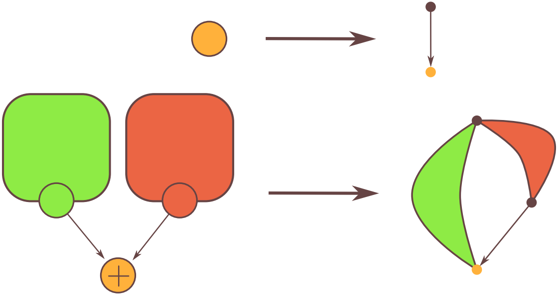

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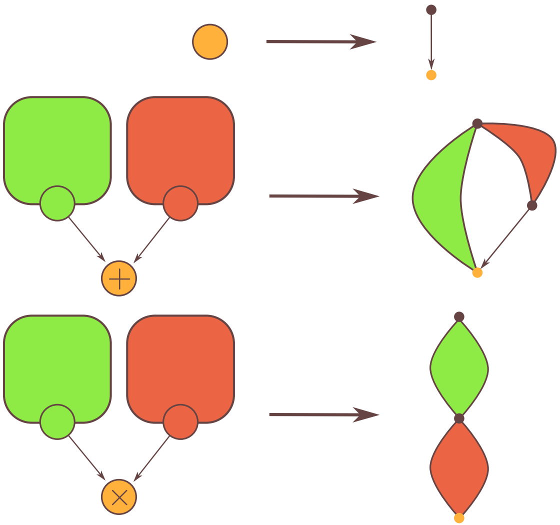

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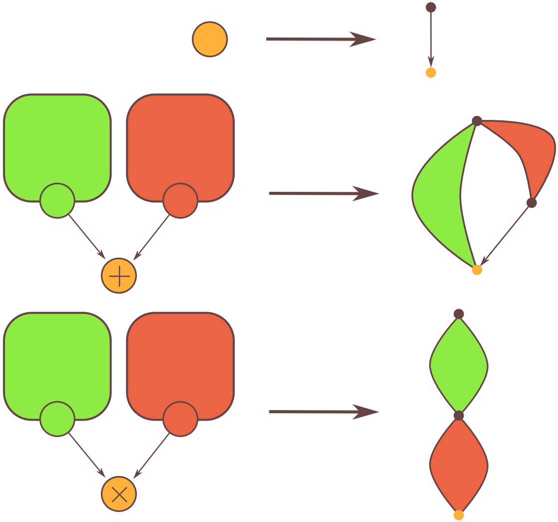

Invariant

$$\varphi = \pm \sum_{s\text{-}t\text{-paths } P} (-1)^{|P|} w(P)$$

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Theorem

For a size-e formula, this construction yields a size-$(e+1)$ graph. Let A be the adjacency matrix of G. Then $\det(A) = \varphi$.

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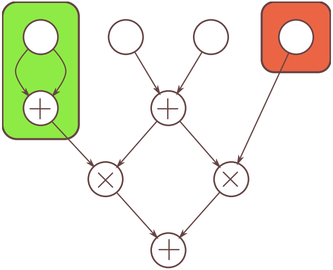

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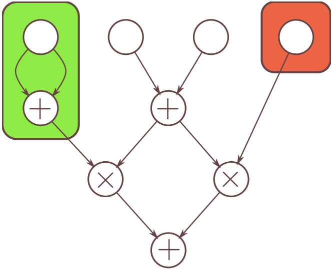

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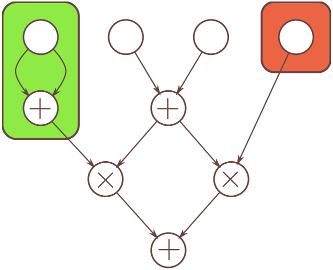

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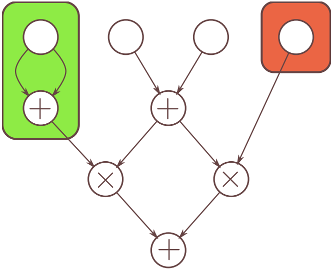

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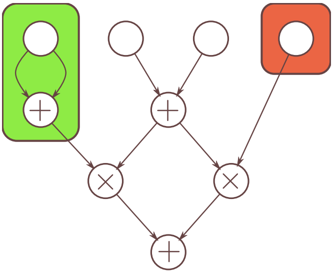

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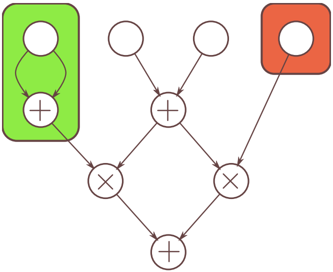

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- Reusable gate: not in a closed subcircuit

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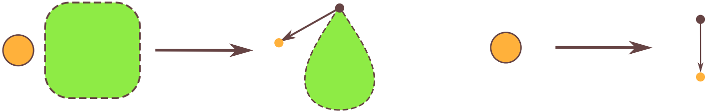

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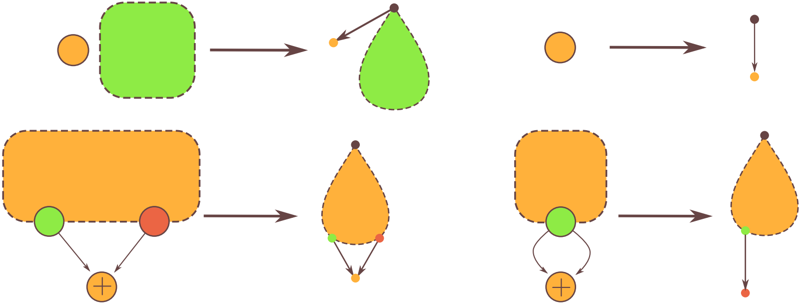

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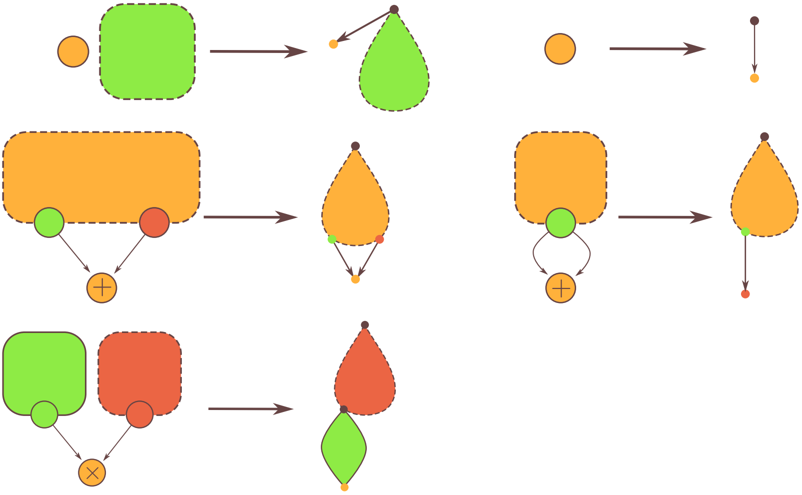

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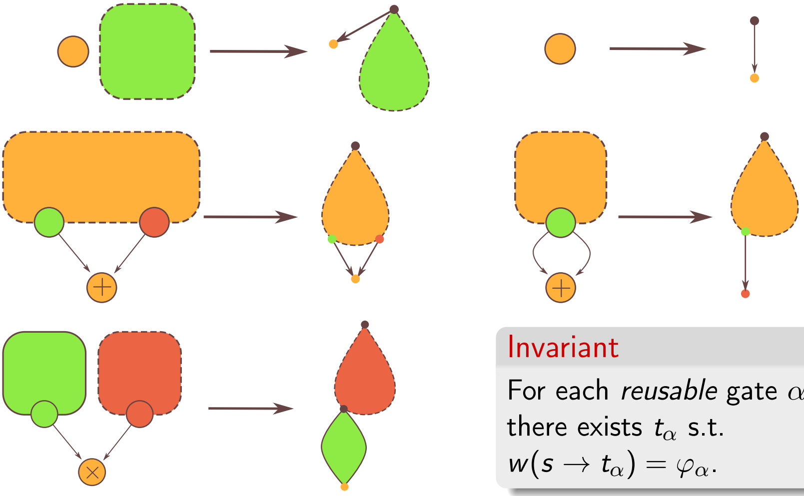

Invariant

For each *reusable* gate α, there exists t_α s.t. $w(s \to t_\alpha) = \varphi_\alpha$.

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Theorem

For a ws circuit of size e with i variable inputs representing φ, this construction yields a size-$(e+i+1)$. The determinant of its adjacency matrix equals φ.

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- Linear Matrix Expression (LME): for A_i symmetric in $\mathbb{R}^{t\times t}$

$$A_0 + x_1 A_1 + \cdots + x_n A_n$$

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- What about polynomial size matrices?

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- N.B.: $\mathrm{char}(\mathbb{K}) \neq 2$ in this section

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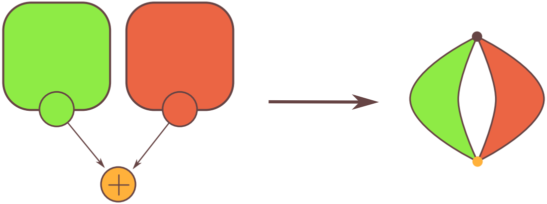

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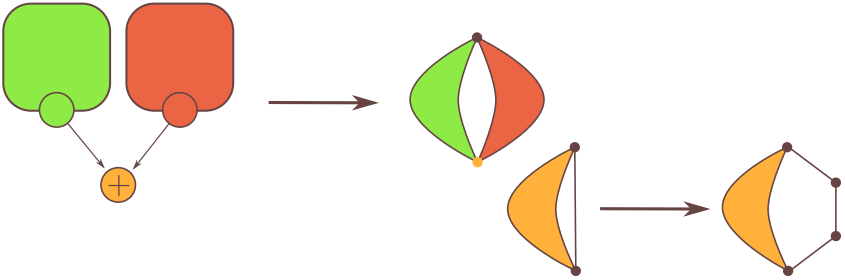

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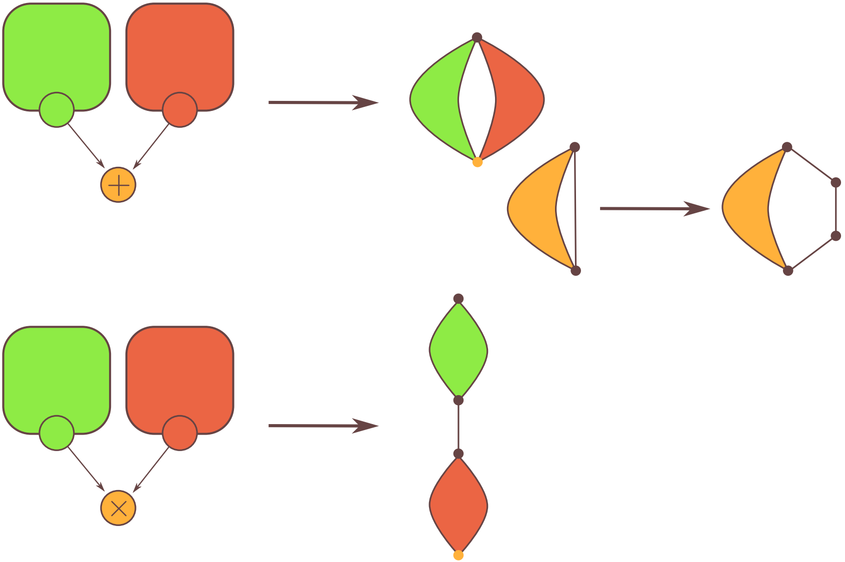

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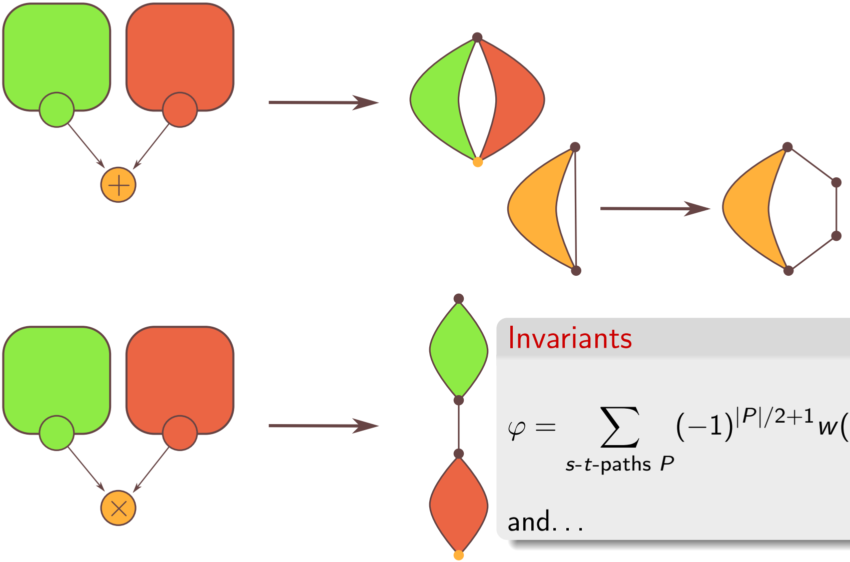

Invariants

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Invariants for formula's construction

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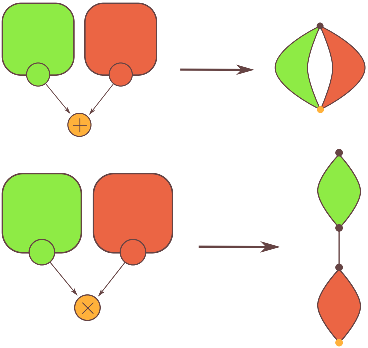

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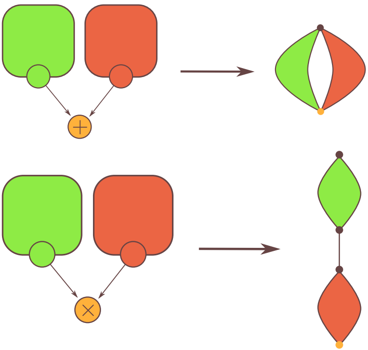

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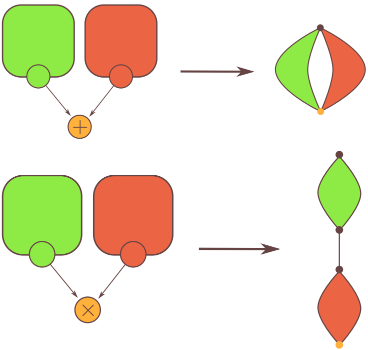

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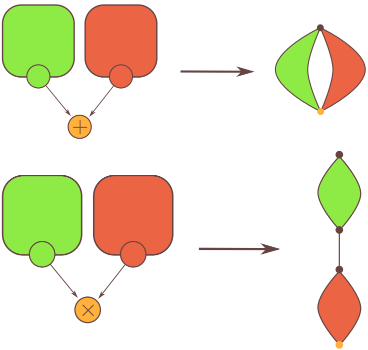

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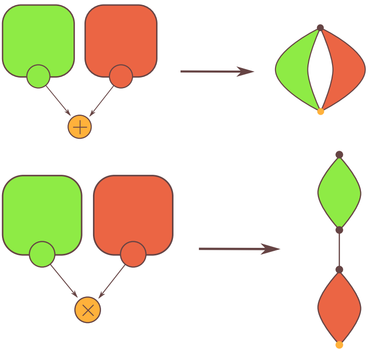

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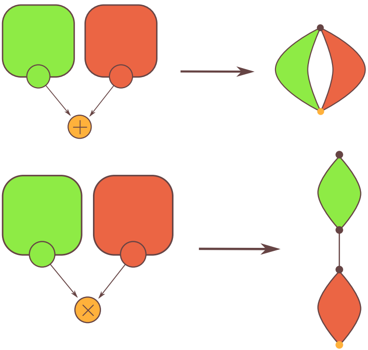

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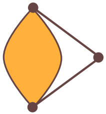

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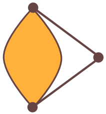

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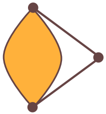

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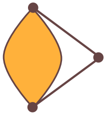

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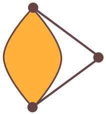

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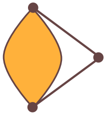

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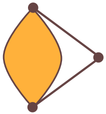

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Theorem

For a formula φ of size e, this construction yields a graph of size $2e+3$. The determinant of its adjacency matrix equals φ.

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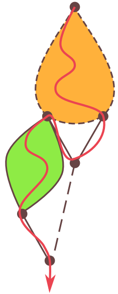

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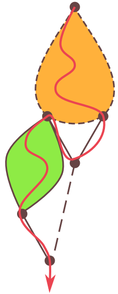

- Definition: an path P is said acceptable if $G \setminus P$ admits a cycle cover

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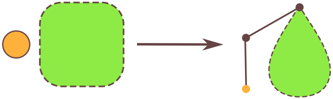

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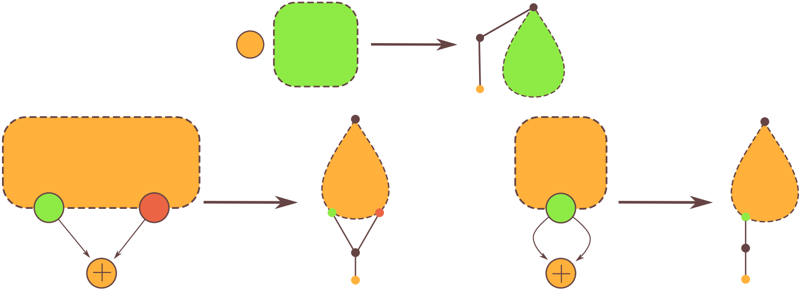

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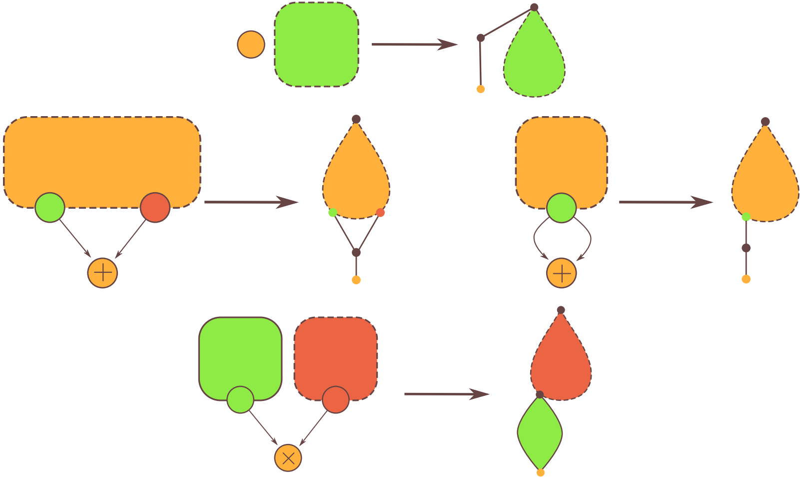

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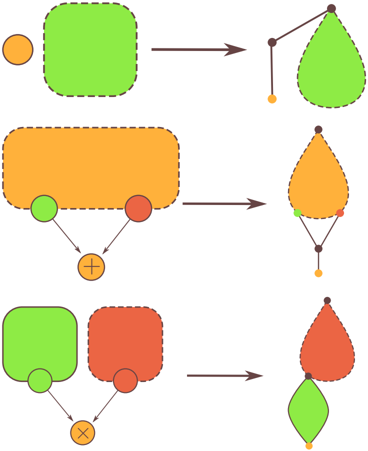

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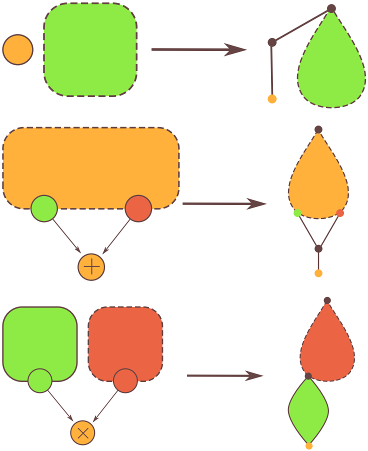

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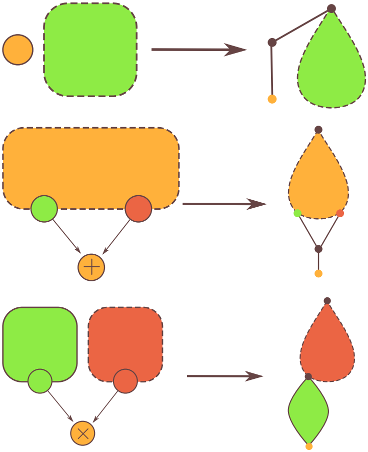

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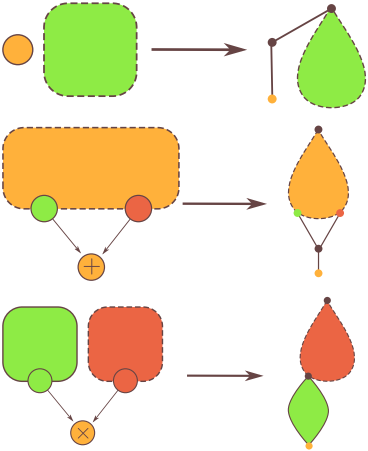

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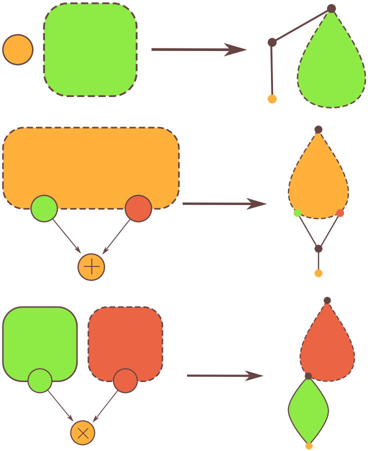

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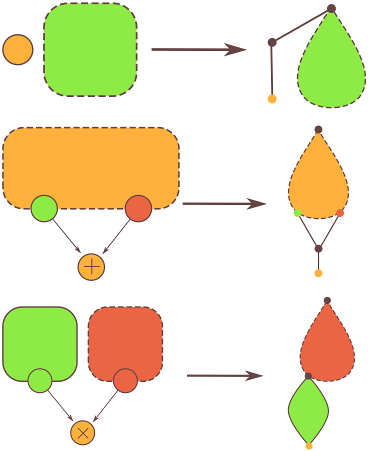

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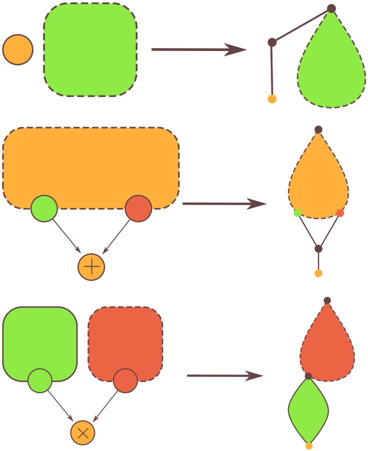

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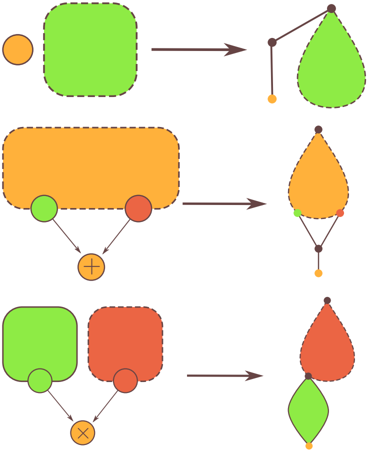

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For a weakly skew circuit of size e, with i input variables, computing a polynomial φ, this construction yields a graph G' with $2(e+i)+1$ vertices. The adjacency matrix of G' has its determinant equal to φ.

Outline

1. From polynomials to determinants
2. From polynomials to determinants of symmetric matrices
3. Characteristic 2

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Which polynomials can be represented as determinant of symmetric matrices in characteristic 2?

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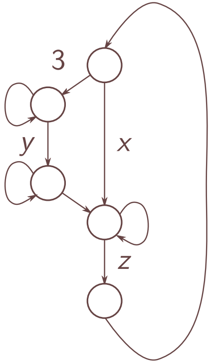

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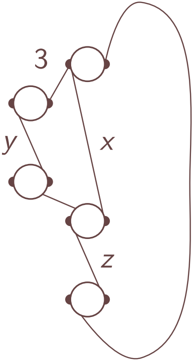

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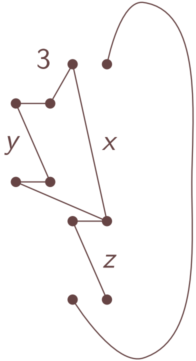

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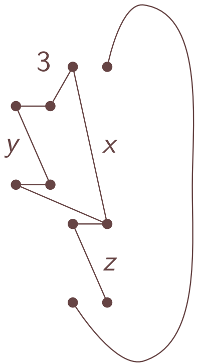

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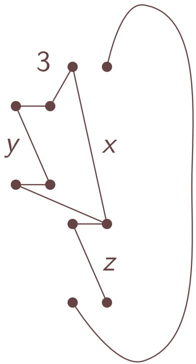

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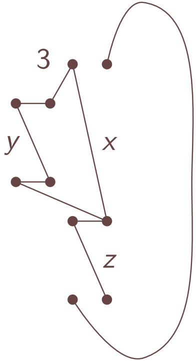

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Is the partial permanent VNP-complete in characteristic 2?

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- $(\mathrm{DET}_n) \in \mathrm{VP}$, $(\mathrm{PER}_n) \in \mathrm{VNP}$, ...

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Same kind of ideas as the previous proof.

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- Symmetric matrices in Valiant's theory?

Thank you!