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**Tutorial 09 – Bool Lyonnais(e)**


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**Exercise 1.**

For all the questions, give the size and the depth of the circuits you build.

1. Give a circuit computing the XOR of its two inputs.
2. Build a circuit computing  $\text{SEL}(x, y, z)$  defined by  $\text{SEL}(0, y, z) = y$  and  $\text{SEL}(1, y, z) = z$ .
3. Let  $f : \{0, 1\}^n \rightarrow \{0, 1\}^m$ . We consider multiple-output circuits. (Equivalently, suppose we are given a circuit for each bit of output.) Show how to turn a multiple-output circuit computing  $f$  into a (classical) circuit deciding  $\{(x, y) : y = f(x)\}$ .
4. Build a multiple-output circuit computing the addition of two binary integers  $x = \overline{x_{n-1} \cdots x_0}$  and  $y = \overline{y_{n-1} \cdots y_0}$  using the schoolbook algorithm.
5. Suppose that  $n$  is a power of 2. Give a circuit computing  $x + y$  using the following recursive algorithm: Compute  $\overline{x_{n-1} \cdots x_{n/2}} + \overline{y_{n-1} \cdots y_{n/2}}$  and  $\overline{x_{n-1} \cdots x_{n/2}} + \overline{y_{n-1} \cdots y_{n/2}} + 1$  in parallel, and use the carry-out of  $\overline{x_{n/2-1} \cdots x_0} + \overline{y_{n/2-1} \cdots y_0}$  to choose the correct answer.
6. Compute the size and depth of the circuit we obtain for  $\{(x, y, z) : x + y = z\}$ . Propose another algorithm to decide this language. Is it better?

**Exercise 2.***Shannon's Effect*

1. Show that every function from  $\{0, 1\}^n$  to  $\{0, 1\}$  can be represented by a CNF formula of size  $n2^n$ .
2. Show that a size- $O(2^n)$  circuit is sufficient.
3. (harder) Are you able to build a size- $O(2^n/n)$  circuit?
4. Show that there exist functions without size- $2^n/(10n)$  circuit.

**Exercise 3.***Spira's Theorem*

A circuit is a *boolean formula* if every gate but the output in it has outdegree exactly 1. That is, the underlying graph is a tree.

- ✎ Show that for every size- $t$  formula  $F$ , there exists an equivalent formula  $F'$  with depth at most  $4 \log t$ . What is the size of  $F'$ ? Can we obtain better bound if  $F'$  is a general circuit instead of a formula?

**Exercise 4.***Monotony*

A boolean function  $f : \{0, 1\}^n \mapsto \{0, 1\}$  is said monotonic if for every  $i$  and every  $x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n$ ,

$$f(x_1, \dots, x_{i-1}, 1, x_{i+1}, \dots, x_n) \geq f(x_1, \dots, x_{i-1}, 0, x_{i+1}, \dots, x_n).$$

- ✎ Show that a function is monotonic iff there exists a circuit with gates AND, OR, 0 and 1, but without NOT gate.